



Norwegian University of
Science and Technology

ABUSE-RESISTANT LOCATION TRACKING: BALANCING PRIVACY AND SAFETY IN THE OFFLINE FINDING ECOSYSTEM

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Introduction

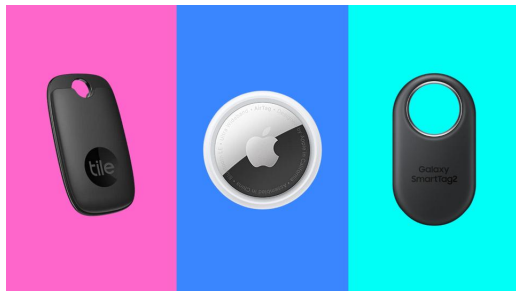
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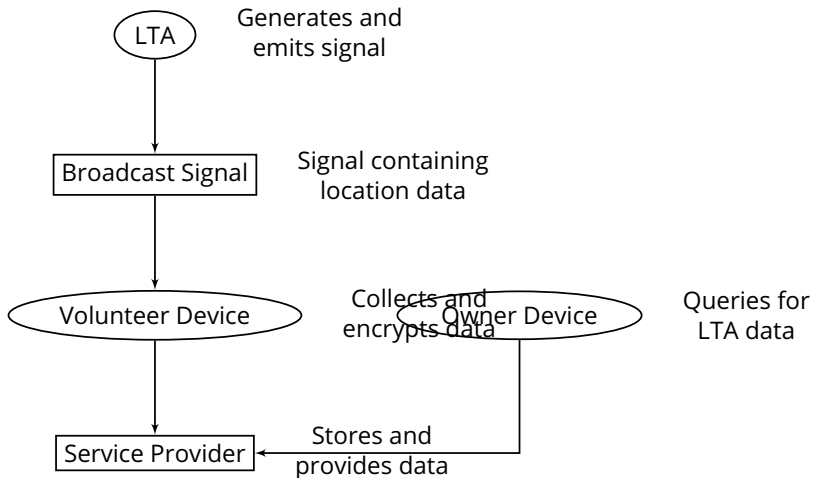
Comparison

Privacy and Security in Location Tracking Accessories (LTAs)



Picture from: <https://www.forbes.com/sites/forbes-personal-shopper/2023/11/30/best-luggage-trackers/>

LTAs



What can they be used for?

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- ▶ Tracking your kids and pets

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- ▶ Tracking your kids and pets
- ▶ Tracking your vehicle
- ▶ Tracking PEOPLE! :(

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- ▶ With tracking comes stalking (unfortunately)
- ▶ The challenge is therefore to design a system that both tracks and can be detected if stalking.
- ▶ This is hard to do because you don't want other people tracking your things either.

Contributions

- ▶ Development of Abuse-Resistant Private Offline Finding Protocols

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- ▶ Experimental Data Collection for Parameter Selection
- ▶ Proposing Protocol Optimizations
- ▶ Implementation and Efficiency Demonstration

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Current State of LTAs

- ▶ Rotate identifiers for privacy

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Current State of LTAs

- ▶ Rotate identifiers for privacy
- ▶ Use secret keys to generate sequences of identifiers
- ▶ Stalker detection becomes harder because of rotating identifiers
- ▶ Attempts to balance privacy and stalker detection
- ▶ Limited against sophisticated tracking

Balancing User Privacy and Stalker Detection

Apple FindMy

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Apple FindMy

- ▶ Pseudonym identifiers

Balancing User Privacy and Stalker Detection

Apple FindMy

- ▶ Pseudonym identifiers
- ▶ Secret key mechanism generating IDs

Balancing User Privacy and Stalker Detection

The Paper

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- ▶ Using MDDS and list-decodable error-correcting codes to enhance privacy

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- ▶ Better stalking detection

Balancing User Privacy and Stalker Detection

The Paper

- ▶ Using MDDS and list-decodable error-correcting codes to enhance privacy
- ▶ Better stalking detection
- ▶ Requires no more trust than Apple FindMy

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Definitions

Definition (Correctness)

A privacy preserving tracking protocol satisfies correctness if for all authorized owners Owner and compliant service providers SP, $\forall loc, aux$ and allowed anonymity epochs i_{epoch} , and $\forall D$ provided by SP.

Definitions

Definition (Detectability)

A privacy preserving tracking protocol is detectable for predicate P' if \forall valid cfg values, $\forall n.u.p.p.t$ algorithms \mathcal{A} , \exists a negligible function $\text{negl}(\lambda)$ so that $\Pr[\text{Exp}_{\mathcal{A}}^{\text{Det}, P'}(\lambda, \text{cfg}) = 0] \leq \text{negl}(\lambda)$.

```
Tag Detectability Experiment

 $Q := \emptyset; L := []; B_{\text{list}} := []; \text{st} = \perp$ 
 $\forall j \in [1, o = \text{poly}(\lambda)] :$ 
   $(\text{id}, i, \text{aux}, \text{loc}, \text{st}) \leftarrow \mathcal{A}(\text{st}, \text{"query"})$ 
  If  $(\text{id}, *) \notin L :$ 
     $k_{\text{tag}} \leftarrow \text{KeyGen}(1^\lambda, \text{cfg})$ 
     $L := L \cup \{(\text{id}, k_{\text{tag}})\}$ 
    find  $k_{\text{tag}}$  so that  $(\text{id}, k_{\text{tag}}) \in L$ 
     $B \leftarrow \text{Beacon}(k_{\text{tag}}, i, \text{aux})$ 
     $B_{\text{list}} := B_{\text{list}} \cup \{(B, \text{loc})\}$ 
     $Q := Q \cup \{(\text{id}, i)\}$ 
out  $\leftarrow \text{Detect}(B_{\text{list}})$ 
 $b = 1$ 
 $\forall (\text{id}, k_{\text{tag}}) \in L,$ 
  If  $P'(\text{cfg}, Q, \text{id}) = 1,$ 
     $b' = (\exists i, (\text{id}, i) \in Q \wedge$ 
       $\text{GetTagID}(k_{\text{tag}}, i) \in \text{out})$ 
     $b = b \wedge b'$ 
return  $b$ 
```

Figure 6: Experiment $\text{Exp}_{\mathcal{A}}^{\text{Det}, P'}(\lambda, \text{cfg})$.

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- ▶ Extension of traditional secret sharing
- ▶ MDSS extends secret sharing by allowing multiple dealers, each with their own secrets, to distribute shares independently.
- ▶ Unlinkability - It's hard to determine if two shares comes from the same dealer or not.
- ▶ Multi-Dealer correctness - Secrets can be accurately reconstructed even with multiple dealers

List-Decodable Error-Correcting Codes

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Application in LTAs: Enables stalker-detection algorithms to function effectively despite identifier rotation.

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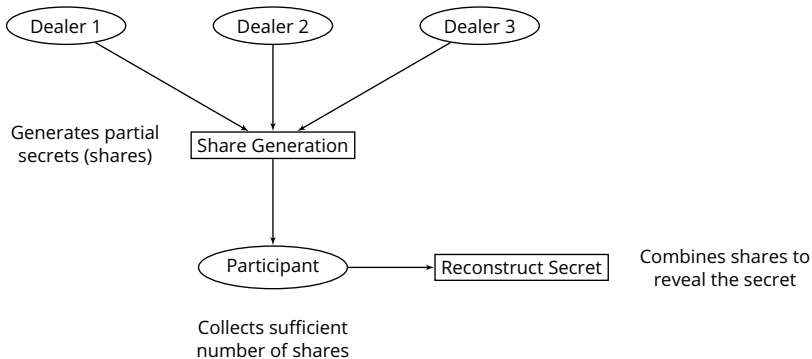
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Application in LTAs: Enables stalker-detection algorithms to function effectively despite identifier rotation.

Benefit: Improves the balance between user privacy (through identifier rotation) and effective stalker detection.

Multi Dealer Secret Sharing (MDSS)



MDSS Construction

Share(s, I')

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Reconstruct($\{sh_1, \dots, sh_{max}\}$)

1. Return $CH^* - MDSS(\{sh_1, \dots, sh_w\})$

CH* Construction

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 1. Uses interpolation to construct polynomials from input
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 3. Reduces the basis
 4. Solution is based on the shortest vector

Algorithm 1: CH*

Input : $\lambda, k, n, \{(\alpha_i, \beta_{i,1}, \dots, \beta_{i,c})\}_{i=1}^n$
Output: None, Multiple, or a single solution $(p_1 \dots p_c)$

- 1 **forall** $j \in [c]$ **do**
- 2 | $f_j(z) = \text{LagrangeInterpol}(\{(\alpha_1, \beta_{1,j}), \dots, (\alpha_n, \beta_{n,j})\})$
- 3 **end**
- 4 $N(z) = \prod_{i=1}^n (z - \alpha_i)$
- 5 construct the matrix $M \in \mathbb{F}[z]^{(c+1) \times (c+1)}$;
- 6

$$M = \begin{bmatrix} z^k & f_1(z) & f_2(z) & \dots & f_c(z) \\ & N(z) & & & \\ & & N(z) & & \\ & & & \ddots & \\ & & & & N(z) \end{bmatrix}$$

- 7 $M_{\text{red}} \leftarrow \text{LatticeReduce}(M)$
- 8 **if** *there is one shortest vector* $\vec{v} = (v_0 \dots v_c)$ *that has a length* $\leq \lambda$ **then**
- 9 | **return** $(v_1 \cdot z^k / v_0, \dots, v_c \cdot z^k / v_0)$
- 10 **end**
- 11 **if** *there are multiple vectors with the shortest length and the shortest length* $\leq \lambda$ **then**
- 12 | **return** Multiple
- 13 **else**
- 14 | **return** None
- 15 **end**

CH*-MDSS Construction

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 1. Applies CH* iteratively
 2. If multiple solutions are found, modify ws to isolate a single solution
 3. Does this with random deletion of coordinates

Algorithm 2: CH*-MDSS Construction

```
Input :  $k, t, n, \{(\alpha_i, \beta_{i,1}, \dots, \beta_{i,c})\}_{i=1}^n$ 
Output: a list  $\{(p_1^i \dots p_c^i)\}_{i=1}^z$  or  $\perp$ 
1 solns := [], ws := [n], fail := False
2 while |ws|  $\geq t$  and not fail do
3   res  $\leftarrow$  CH*( $k + (|ws| - t)$ ,  $k, n, \{(\alpha_i, \beta_{i,1} \dots \beta_{i,c})\}_{i \in ws}$ ).
4   if res = None then
5     | return solns
6   else if res = Multiple then
7     |  $s \leftarrow \lfloor \frac{|ws|}{t} \rfloor$ , found  $\leftarrow$  False
8     | while not found: do
9       | // randomly remove points until a single solution is found
10      | pts  $\leftarrow$  ws.pop( $2(s - 1)$ )
11      | res  $\leftarrow$  CH*( $k + (|ws| - t)$ ,  $k, n, \{(\alpha_i, \beta_{i,1} \dots \beta_{i,c})\}_{i \in ws}$ )
12      | ws  $\leftarrow$  ws  $\cup$  pts
13      | if res can be parsed as  $(p_1 \dots p_c)$  then
14        | found  $\leftarrow$  True
15        | add  $(p_1 \dots p_c)$  to solns, remove agreeing points from ws
16      | end
17    | endw
18  | else
19  | parse res as  $(p_1 \dots p_c)$ , add  $(p_1 \dots p_c)$  to solns, remove agreeing points from ws
20  | end
21 endw
return solns
```

Abuse-resistant Offline Finding Protocol

$\text{KeyGen}(1^\lambda, \text{cfg}) \rightarrow k_{\text{tag}}$ In: required params, out: a secret key.

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report.

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$\text{GetTagID}(i_{\text{tag}}, i_{\text{epoch}}) \rightarrow \text{id}_{\text{tag}}$ Helper algorithm used by the LTA and owner device to find the current identifier of the LTA

$\text{GenReport}(B, \text{loc}) \rightarrow R$ In: A broadcast and location, out: a report.

$\text{Detect}(\text{cfg}, \{(B_1, \text{loc}_1), \dots, (B_n, \text{loc}_n)\}) \rightarrow \{\text{id}_{\text{tag}_i}\}$ In: set of broadcasts, out: one or more identifiers.

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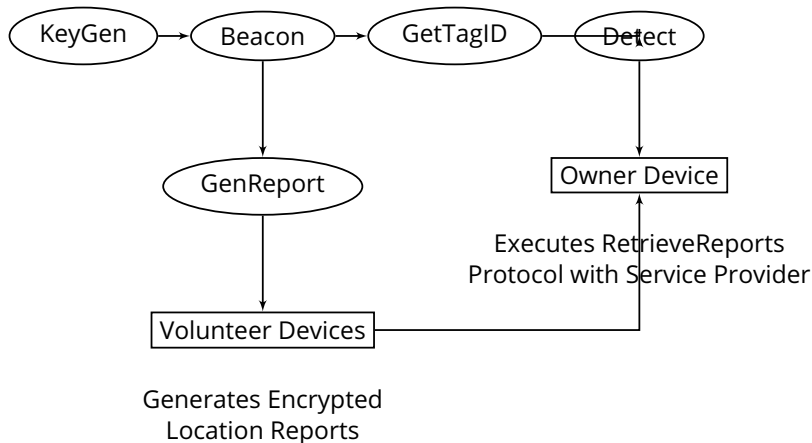
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$\text{RetrieveReports}(\text{Owner}(k_{\text{tag}}, i_{\text{epoch}}), \text{SP}(\mathcal{D}))$ Executed between 2 parties, Owner provides a tag key and epoch, SP uses a database \mathcal{D} , outputting a list of reports.

Offline Finding Protocol Flow



Main Construction: $\text{KeyGen}(\text{cfg})$

1. $k_1 \leftarrow \text{PRF}_1.\text{KeyGen}(1^\lambda)$

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3. $k_3 \leftarrow \text{PRF}_3.\text{KeyGen}(1^\lambda)$
4. return $(k_1, k_2, k_3, \text{cfg})$

Main Construction: Beacon($k_{\text{tag}}, i_{\text{epoch}}, \text{aux}$)

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1. $(k_1, k_2, k_3, \text{cfg}) \leftarrow k_{\text{tag}}$
2. $(E, L) \leftarrow \text{cfg}$
3. $(\text{pk}, _) \leftarrow \text{CCA.KeyGen}(1^\lambda; \text{PRF.Eval}(k_1, i))$

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6. $I := \{0, \dots, L - 1\}$
7. $\text{sh}_0^e \dots \text{sh}_{L-1}^e \leftarrow \text{Share}(\text{id}_{\text{tag}}, I; \text{PRG}(\text{PRF}_2.\text{Eval}(k_2, e)))$

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8. Return $\text{pk} \parallel \text{sh}_i^e \parallel \text{aux}$

Main Construction: $\text{Detect}(\text{cfg}, \{B_j\})$

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2. Return $\text{II.Reconstruct}(S)$

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- ▶ Unlinkability and threshold privacy from MDSS

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Achieving Privacy Goals

- ▶ Unlinkability and threshold privacy from MDSS
- ▶ Ensure identifiers change periodically
- ▶ Should need a copy of the secret key to re-derive the correct sequence to obtain location reports

Stalker Detection

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- ▶ These secret shares are only used locally for stalker detection.

Contents

Introduction

Background

Protocol Explanation

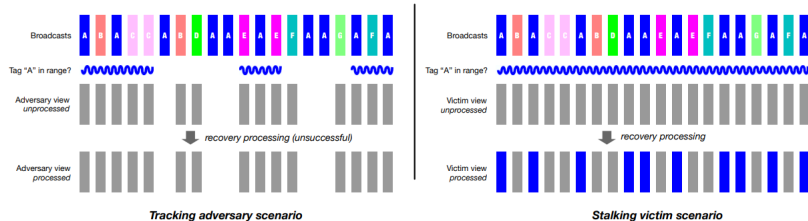
Security Analysis

Comparison

FindMy vs. This Construction

Protocol	Epoch duration	Broadcasts per epoch	Stalker detection?	Tracking privacy	Continuous Proximity	Stalker detection
<i>Apple FindMy [2] / IETF [32]:</i>						
Near-owner mode	15 min	450	✗	n/a	n/a	
Separated mode	24 hrs	43,200	●	n/a	✗	15-60 min [†]
<i>This work (§4):</i>						
2-second epochs / 1-hour window	2 sec	1	●	40 – 46 min*	●	60 min
4-second epochs / 1-hour window	4 sec	1	●	39 – 46 min*	●	60 min
1-minute epochs / 1-hour window	60 sec	15	●	41 – 47 min*	●	60 min

2 Detection Scenarios



Questions?